



The Invention of Scheme, 1762 (colorized)

# Hi! Welcome to 61A Discussion :)

We will begin at **8:10!**

Attendance form and skeleton notes:

**[cs61a.bencuan.me](https://cs61a.bencuan.me)**



**Secret word:**

# Announcements

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- ▣ HW6 due tonight
- ▣ Scheme Part 1 due next Tues.
  - ▣ Start early!

# Agenda

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- ▣ Attendance
- ▣ Tail Recursion
- ▣

# Tail Recursion

# The Vocab

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- ❑ **Tail Context:** the **very last** thing that's done in a Scheme expression
- ❑ **Tail Call:** calling a function in tail context
- ❑ **Tail Recursion:** recursive tail call (in tail context)

# Why do we care about tail recursion?

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Consider this implementation of `factorial` that is *not* tail recursive:

```
(define (factorial n)
  (if (= n 0)
      1
      (* n (factorial (- n 1)))))
```

Here's a visualization of the recursive process for computing `(factorial 6)` :

```
(factorial 6)
(* 6 (factorial 5))
(* 6 (* 5 (factorial 4)))
(* 6 (* 5 (* 4 (factorial 3))))
(* 6 (* 5 (* 4 (* 3 (factorial 2))))))
(* 6 (* 5 (* 4 (* 3 (* 2 (factorial 1)))))))
(* 6 (* 5 (* 4 (* 3 (* 2 1))))))
(* 6 (* 5 (* 4 (* 3 2))))
(* 6 (* 5 (* 4 6)))
(* 6 (* 5 24))
(* 6 120)
720
```

# Factorial, tail recursive version

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```
(define (factorial n)
  (define (fact-tail n result)
    (if (= n 0)
        result
        (fact-tail (- n 1) (* n result))))
  (fact-tail n 1))
```

# Scheme project EC ideas

- ▣ Instead of creating a new frame, replace the old frame (since you don't need it anymore)





## Q1: Is Tail Call

For each of the following procedures, identify whether it contains a recursive call in a tail context. Also indicate if it uses a constant number of active frames.

```
(define (question-a x)
  (if (= x 0) 0
      (+ x (question-a (- x 1)))))
```

```
(define (question-b x y)
  (if (= x 0) y
      (question-b (- x 1) (+ y x))))
```

### ▣ **Strategy:**

1. What is the last operation made?
2. Is it a recursive call?
3. Is it normal recursion or tree recursion?

## Q1: Is Tail Call

For each of the following procedures, identify whether it contains a recursive call in a tail context. Also indicate if it uses a constant number of active frames.

```
(define (question-c x y)
  (if (> x y)
      (question-c (- y 1) x)
      (question-c (+ x 10) y)))
```

```
(define (question-d n)
  (if (question-d n)
      (question-d (- n 1))
      (question-d (+ n 10))))
```

```
(define (question-e n)
  (cond ((<= n 1) 1)
        ((question-e (- n 1)) (question-e (- n 2)))
        (else (begin (print 2) (question-e (- n 3))))))
```

### ▣ Strategy:

1. What is the last operation made?
2. Is it a recursive call?
3. Is it normal recursion or tree recursion?

# Writing tail recursive functions

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- 1. Create a helper function!**
  - a. Parameters: so-far, anything else that changes
- 2. Base case:** return so-far
- 3. Recursive case:** must be in tail context
  - a. Change so-far
  - b. Very first operator on the line should be function name
- 4. Remember to call the helper function!**

## Q2: Sum

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Write a tail recursive function that takes in a Scheme list and returns the numerical sum of all values in the list. You can assume that the list contains only numbers (no nested lists).

```
scm> (sum '(1 2 3))  
6  
scm> (sum '(10 -3 4))  
11
```

# Q3: Reverse

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## Q3: Reverse

Write a tail-recursive function `reverse` that takes in a Scheme list and returns a reversed copy. *Hint*: use a helper function!

```
scm> (reverse '(1 2 3))
(3 2 1)
scm> (reverse '(0 9 1 2))
(2 1 9 0)
```

# Interpreters

# The Calculator Example

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- ▣ **Goal:** let's write an interpreter that understands simple math expressions!
  - ▣  $(+ 2 2)$
  - ▣  $(- 5)$
  - ▣  $(* (+ 1 2) (+ 2 3))$
- ▣ Understands  $+$ ,  $-$ ,  $*$ ,  $/$ , and nested expressions

# Pairs

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- ▣ Pairs are literally linked lists!! Only differences:
  - ▣ **Pair** vs Link
  - ▣ **nil** vs Link.empty
  - ▣ **Pair(1, nil)** vs **Link(1)**
- ▣ Used to represent Scheme code in Python



# Operators and Operands

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- ▣ An **operator** is the function you are trying to apply in Scheme
  - ▣ +, list, append...
  - ▣ The **very first** element in a Pair list
- ▣ An **operand** is a parameter that is passed into the function
  - ▣ In (+ 3 5), + is the operator and 3, 5 are both operands

## Q4: Using Pair

Answer the following questions about a `Pair` instance representing the Calculator expression `(+ (- 2 4) 6 8)`.

Write out the Python expression that returns a `Pair` representing the given expression:

What is the operator of the call expression?

If the `Pair` you constructed in the previous part was bound to the name `p`, how would you retrieve the operator?

What are the operands of the call expression?

If the `Pair` you constructed was bound to the name `p`, how would you retrieve a list containing all of the operands?

How would you retrieve only the first operand?

**Hint:** Pairs have `.first` and `.rest` just like a linked list! **do not** use `car`, `cdr`, `cons`...

# Q4 Solutions

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Write out the Python expression that returns a `Pair` representing the given expression:

Draw a box and pointer diagram corresponding to the `Pair`:

# Q4 Solution 2

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What is the operator of the call expression?

If the `Pair` you constructed in the previous part was bound to the name `p`, how would you retrieve the operator?

What are the operands of the call expression?

If the `Pair` you constructed was bound to the name `p`, how would you retrieve a list containing all of the operands?

How would you retrieve only the first operand?

## Q9: From Pair to Calculator

Write out the Calculator expression with proper syntax that corresponds to the following `Pair` constructor calls.

```
>>> Pair('+', Pair(1, Pair(2, Pair(3, Pair(4, nil))))))
```

```
>>> Pair('+', Pair(1, Pair(Pair('*', Pair(2, Pair(3, nil))), nil)))
```

# Eval and Apply

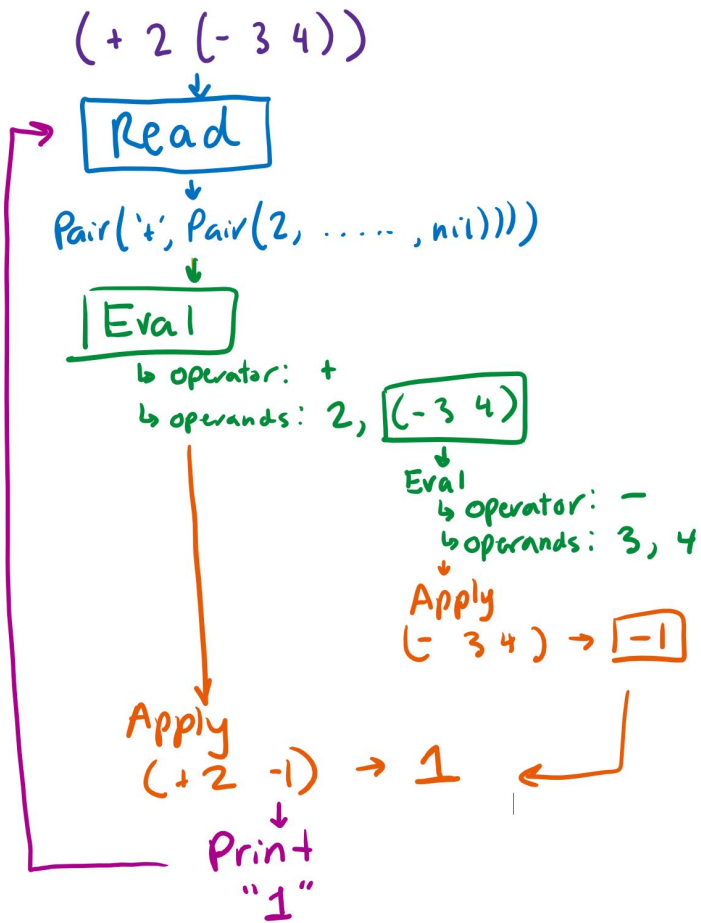
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## ▣ Eval rules:

- ▣ Basic elements (numbers, booleans...) - done (base case)
- ▣ Built-in functions (+, -, ...) - lookup in OPERATORS
- ▣ Function:
  - Call eval on operator
  - Recursively call eval on all operands
  - Call apply on operator and operands

# Example

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# Counting eval and apply

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- ▣ Entire expression = +1 eval
- ▣ Operator = +1 eval, +1 apply
- ▣ Operand = +1 eval
- ▣ **Remember to recursively run eval on nested expressions!**
- ▣ Example:  $(+ 2 (- 3 4)) \rightarrow 7$  eval, 2 apply
- ▣ **Do not count special forms**
  - ▣ Example:  $(\text{and } 1 2 3) \rightarrow 2$  eval, 0 apply



## Q8: Counting Eval and Apply

How many calls to `calc_eval` and `calc_apply` would it take to evaluate each of the following Calculator expressions?

```
scm> (+ 1 2)
```

```
scm> (+ 2 4 6 8)
```

```
scm> (+ 2 (* 4 (- 6 8)))
```

```
scm> (and 1 (+ 1 0) 0)
```

- ▣ **Entire expression = +1 eval**
- ▣ **Operator = +1 eval, +1 apply**
- ▣ **Operand = +1 eval**
- ▣ **Remember to recursively run eval on nested expressions!**
- ▣ Example: `(+ 2 (- 3 4))` → 7 eval, 2 apply
- ▣ **Do not count special forms**
- ▣ Example: `(and 1 2 3)` → 2 eval, 0 apply

## Q8: Counting Eval and Apply

How many calls to `calc_eval` and `calc_apply` would it take to evaluate each of the following Calculator expressions?

```
scm> (+ 1 2)
```

```
scm> (+ 2 4 6 8)
```

## Q8: Counting Eval and Apply

How many calls to `calc_eval` and `calc_apply` would it take to evaluate each of the following Calculator expressions?

```
scm> (+ 2 (* 4 (- 6 8)))
```

```
scm> (and 1 (+ 1 0) 0)
```

# New Procedure! (Q5)

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Suppose we want to add the `//` operation to our Calculator interpreter. Recall from Python that `//` is the floor division operation, so we are looking to add a built-in procedure `//` in our interpreter such that `(// dividend divisor)` returns `dividend // divisor`. Similarly we handle multiple inputs as illustrated in the following example `(// dividend divisor1 divisor2 divisor3)` evaluates to `((dividend // divisor1) // divisor2) // divisor3`. For this problem you can assume you are always given at least 1 divisor. Also for this question do you need to call `calc_eval` inside `floor_div`? Why or why not?

# New Procedure! (Q5)

```
1 def calc_eval(exp):
2     if isinstance(exp, Pair): # Call expressions
3         return calc_apply(calc_eval(exp.first), exp.rest.map(calc_eval))
4     elif exp in OPERATORS:   # Names
5         return OPERATORS[exp]
6     else:                    # Numbers
7         return exp
8
9 def floor_div(expr):
10     """
11     >>> calc_eval(Pair("/", Pair(10, Pair(10, nil))))
12     1
13     >>> calc_eval(Pair("/", Pair(20, Pair(2, Pair(5, nil)))))
14     2
15     >>> calc_eval(Pair("/", Pair(6, Pair(2, nil))))
16     3
17     """
18     """ YOUR CODE HERE """
19
20 OPERATORS = { "/": floor_div }
21
```

# New Form! (Q6)

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- ▣ Can we add  $\leq$ ,  $\geq$ , etc. without changing `calc_eval`?
- ▣ What about `and`, `or`?

```
def calc_eval(exp):
    """Evaluates a Calculator expression represented as a Pair."""
    if isinstance(exp, Pair): # Call expressions
        fn = calc_eval(exp.first)
        args = list(exp.rest.map(calc_eval))
        return calc_apply(fn, args)
    elif exp in OPERATORS: # Names
        return OPERATORS[exp]
    else: # Numbers
        return exp
```

# New Form! (Q6)

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iii. Now, complete the implementation below to handle `and` expressions. You may assume the conditional operators (e.g. `<`, `>`, `=`, etc) have already been implemented for you.

```
def calc_eval(exp):
    if isinstance(exp, Pair):

        if _____: # and expressions
            return eval_and(exp.rest)
        else: # Call expressions
            return calc_apply(calc_eval(exp.first), list(exp.rest.map(calc_eval)))
    elif exp in OPERATORS: # Names
        return OPERATORS[exp]
    else: # Numbers
        return exp

def eval_and(operands):
    """ YOUR CODE HERE """
```

# Define (Q7)

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In the last few questions we went through a lot of effort to add operations so we can do most arithmetic operations easily. However it's a real shame we can't store these values. So for this question let's implement a `define` special form that saves values to variable names. This should work like variable assignment in Scheme; this means that you should expect inputs of the form `(define <variable_name> <value>)` and these inputs should return the symbol corresponding to the variable name.

```
calc> (define a 1)
a
calc> a
1
```

This is a more involved change. Here are the 4 steps involved:

1. Add a `bindings` dictionary that will store the names and corresponding values of variables as key-value pairs of the dictionary.
2. Identify when the `define` form is given to `calc_eval`.
3. Allow variables to be looked up in `calc_eval`.
4. Write the function `eval_define` which should actually handle adding names and values to the `bindings` dictionary.